

# Kaa: A Python Implementation of Reachable **Set Computation Using Bernstein Polynomials**

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- of dynamical and hybrid systems.
- bundles.
- implementation with only around ~650 lines of code.

#### Introduction

• Reachable set computation is one of the many important tools available for the verification

• One of the simpler and easier-to-understand reachable set computation algorithms for polynomial discrete dynamical systems utilizes Bernstein polynomials and parallelotope

• Tomasso Dreossi, Thao Dang and Carla Piazza implemented a tool called Sapo in C++ which leverages parallelotope bundles and the properties of Bernstein polynomials.

• Kaa is a reimplementation of Sapo using robust Python libraries. The result is a compact







- The state of a system, denoted as x, lies in a domain  $D \subseteq \mathbb{R}^n$ . A discrete-time polynomial nonlinear system is denoted as  $x^+ = f(x)$
- The trajectory is denoted as  $\xi(x_0)$ , is the sequence  $x_0, x_1, \dots$  where  $x_{i+1} = f(x_i)$ .
- Given an initial set  $\Theta$ , the reachable set at time k, denoted as  $\Theta_k = \{ \xi(x_0, k) \mid x_0 \in \Theta \}$  where  $\xi(x_0, k) = x_k$ .

#### Preliminaries



- A parallelotope P is a set of states in  $\mathbb{R}^n$  denoted as  $\langle \Lambda, c \rangle$  where  $\Lambda \in \mathbb{R}^{2n \times n}$  and  $c \in \mathbb{R}^{2n}, \Lambda_{i+n} = -\Lambda_i$  and  $i \in \{1, \dots, n\}$  such that:  $x \in P$  if and only if  $\Lambda x \leq c$ .
- A is called the *direction matrix* where  $\Lambda_i$  denotes the  $i^{th}$ row of  $\Lambda$ . The vector c is called the offset vector where  $c_i$  is the  $i^{th}$  element of the vector.
- A parallelotope bundle Q is a set of parallelotopes  $\{P_1, \ldots, P_m\}$  where  $Q = \bigcap_{i=1}^m P_i$ . Note that any polytope initial set can be expressed as a parallelotope bundle.

### Parallelotope Bundles









### Parallelotope Bundles

Figure 1 from Dreossi et. al: Parallelotope Bundles for Polynomial Reachability (2016)





degree d and index i is

 $\mathscr{B}_{i,d} = \beta_{i_1,d_1}(x_1)\beta_{i_2,d_2}(x_2)\dots\beta_{i_n,d_n}(x_n)$ 

 $\beta_{i_m,d_m}(x_m) = \begin{pmatrix} d_m \\ i_m \end{pmatrix} x_m^{i_m}(1-x_m)^{d_m-i_m}$ 

• Any polynomial function can be expressed in the Bernstein basis.

#### Bernstein Polynomials

• Given two multi-indices i and d of size n, where  $i \leq d$ , the Bernstein polynomial of







• The corresponding *Bernstein Coefficients* can be explicitly calculated for multiindex *i* and polynomial degree *d*:

$$b_{i,d} = \sum_{j \le i} \prod_{r} \frac{\binom{i_r}{j_r}}{\binom{d_r}{j_r}} a_j$$

The upper and lower bounds of polynomial  $h(x_1, ..., x_n)$  over unit box  $[0,1]^n$  are bounded by the Bernstein coefficients:

$$\min_{i\in I}\{b_i\} \leq \inf_{x\in[0,1]^n}h(x)$$

### **Bernstein Polynomials**

$$\leq \sup_{x \in [0,1]^n} h(x) \leq \max_{i \in I} \{b_i\}.$$





- to P.
- upper bound of  $h \circ T_p$  over  $[0,1]^n$ .
- polynomial h over some parallelotope P as BernsteinUpper(h, P) and BernsteinLower(h, P) respectively.

### Reachable Set Comp.

• A parallelotope P can also be represented as an affine transformation  $T_p$  from  $[0,1]^n$ 

• Therefore, upper bounds on the supremum of a function h over P is equivalent to

We denote the procedures for calculating such upper and lower bounds for a







- we wish to compute an over-approximation of the image f(Q) as a new bundle  $Q' = \{P'_1, P'_2, \dots, P'_m\}.$
- problems:

 $c_{j,i} = m_{x \in x}$ 

 $c_{j+n,i} = \max_{x \in x}$ 

matrix for  $P_i$ .

### **Reachable Set Comp.**

• Given parallelotope bundle  $Q = \{P_1, P_2, \dots, P_m\}$  and a discrete dynamical system  $x^+ = f(x)$ ,

• We ensure that direction matrix  $\Lambda_{i}$  of  $P'_{i}$  is same as  $P_{i}$  and the computation is required only to compute the offsets of the directions according to the following non-linear optimization

$$\begin{aligned} & \underset{i \in P_i}{\max} \Lambda_{j,i} \cdot f(x) \\ & \underset{i \in P_i}{\max} - \Lambda_{j,i} \cdot f(x) \end{aligned}$$

• Here,  $c_{j,i}$  is the  $j^{th}$  offset of parallelotope  $P_i$ . Similarly,  $\Lambda_{j,i}$  is the  $j^{th}$  row of the directions









 $Q = \{P_1, P_2, \dots, P_m\}$ :

$$c_{j,i} = \min_{l=1}^{m} \left\{ \text{BernsteinUpper} \left( \Lambda_{j,i} \cdot f(x), P_l \right) \right\} \text{ if } j \leq n .$$
  
$$c_{j+n,i} = \max_{l=1}^{m} \left\{ \text{BernsteinLower} \left( \Lambda_{j,i} \cdot f(x), P_l \right) \right\} \text{ otherwise.}$$

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### Reachable Set Comp.

• We can invoke BernsteinUpper(h, P) and BernsteinLower(h, P) to update the offsets according to the solutions found over all parallelotopes in the bundle

We iterate this over a certain number of time steps to produce the reachable set.





- polynomial dynamical systems (2017).
- code.
- papers to find an explanation of the inner workings.
- Consequently, it becomes difficult to accommodate experimentation.

#### Sapo Drawbacks

First published in: Dreossi, T.: Sapo: Reachability computation and parameter synthesis of

Current implementation is verbose. The main core of algorithm takes over a thousand lines of C++

It does not have native plotting functionality. Sapo generates MATLAB code which must be separately run through either MATLAB or Octave. Simultaneous visualization is clunky at best.

Suffers from little to no documentation. The curious reader must delve into previously published







- Python is known for its powerful, well-tested symbolic and matrix-computation libraries.
- *Numpy* libraries are popular matrix-computation libraries which allow higher-level manipulation of matrices. This gives us an avenue of overcoming the verbosity and the possibility of memory leaks inherent in implementing identical features in C++.
- The library of Sympy has powerful symbolic manipulation tools which allow us to comfortably perform many sensitive symbolic subsititutions into polynomials.
- *Matplotlib* library has intuitive plotting facilities that we integrate into our tool for visualizing the reachable set. In particular, *Matplotlib* facilitates the ability to visualize several reachable sets simultaneously.

### Motivations for Kaa







- We offer a Juypter Notebook to rapidly introduce the interested reader to the techniques and tools we offer through Kaa.
- Jupper notebooks are simple to create and well-known for their straightforward user interface.
- By leveraging the *Matplotlib* library for visualizing the reachable set, we were able to design an engaging interactive tutorial and experimentation platform for visualizing reachable sets of non-linear systems.
- We document the code extensively and offer resources for learning the internals of Kaa.

#### Accessibility







• The SIR epidemic model is a 3-dimensional dynamical system governed by the following dynamics:

- $s_{k+1} = s_k (\beta s_k i_k) \Delta$  $\beta = 0.34, \gamma = 0.05$  $i_{k+1} = i_k + (\beta s_k i_k - \gamma i_k)\Delta$  $\Delta = 0.5$  $r_{k+1} = r_k + (\gamma i_k)\Delta$

Time Steps	Kaa (offu)	Sapo (offu)	Kaa (offl)	Kaa (offl)
50	0.470716	0.470716	-0.435191	-0.43519
51	0.475839	0.475839	-0.439599	-0.439599
52	0.480906	0.480906	-0.443945	-0.443945
53	0.485915	0.485915	-0.448227	-0.448227
54	0.490862	0.490862	-0.452443	-0.452443
55	0.495747	0.495747	-0.456591	-0.456591
56	0.500566	0.500566	-0.460669	-0.460669
57	0.505317	0.505317	-0.464675	-0.464675
58	0.509999	0.509999	-0.4686075	-0.468608
59	0.514610	0.514610	-0.472465	-0.472465
60	0.519147	0.519147	-0.476246	-0.476246

Upper and lower offsets for variable I

### **Results: SIR Model**





Kaa



#### **Results: SIR Model**





Kaa



### **Results: SIR Model**

Sapo





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### **Results:** Rossler Model

- $x_{k+1} = x_k + -(y z)\Delta$  $y_{k+1} = y_k + (x_k + ay_k)\Delta$
- $z_{k+1} = z_k + (b + z_k(x_k c))\Delta$

Time Steps	Kaa (offu)	Sapo (offu)	Kaa (offl)	Kaa (offl)
50	1.95908	1.96043	-1.9209	-1.9193
51	1.83552	1.83688	-1.7963	-1.7947
52	1.71044	1.71181	-1.67016	-1.6685
53	1.58392	1.58531	-1.54255	-1.5409
54	1.45604	1.45744	-1.41355	-1.4119
55	1.32687	1.32829	-1.28324	-1.2815
56	1.19649	1.19792	-1.15168	-1.1500
57	1.06499	1.06642	-1.01896	-1.0172
58	0.932432	0.933877	-0.885157	-0.88339
59	0.798905	0.800358	-0.75035	-0.74857
60	0.664489	0.665949	-0.614619	-0.61282

Upper and lower offsets for variable y

• The Rossler model is another 3-dimensional system governed under the dynamics:

a = 0.1, b = 0.1, c = 14 $\Delta = 0.025$ 





Kaa



#### **Results: Rossler Model**

Sapo



#### Kaa



#### **Results: Rossler Model**





#### Kaa



# **Results: Quadcopter Model**

Sapo







- computation.
- improve on them.

Model	Kaa	$\mathbf{SAPO}  (C++)$
SIR	$11.41  \sec$	0.16 sec
Rossler	$41.92  \sec$	$1.17  \mathrm{sec}$
Quadcopter	$78.21  \sec$	$11.98  \sec$
Lotka-Volterra	$18 \min 95.05 \sec$	$57.48  \sec$
Phosphoraley	$103.81  \sec$	24.86 sec

#### Performance Drawbacks

While the current implementation in Python is very intuitive and concise, it incurs severe performance penalties. We believe this is due to some extraneous library calls in the core loop of the reachable set

An immediate next step is to deploy extensive profiling to find performance bottlenecks and subsequently







- which is focused towards accessibility and pedagogical use.
- Formal-Methods/kaa
- top of the library.
- streamlined format for defining models and visualizing reachable sets.

#### Conclusions

We present Kaa, a Python implementation of reachable set computation of nonlinear systems

We include Juypter Notebooks and documentation through: <u>https://github.com/Tarheel-</u>

While we do incur performance drawbacks from selecting Python for implementing this algorithm, we believe that it aids in fast prototyping and enables students to easily build on

Immediate future work includes improving on the running time and creating a more







- Dreossi, T.: Sapo: Reachability computation and parameter synthesis of polynomial dynamical systems. In: Proceedings of the 20th International Conference on Hybrid Systems: Computationand Control. pp. 29–34 (2017)
- Dreossi, T., Dang, T., Piazza, C.: Parallelotope bundles for polynomial reachability. In: Proceedings of the 19th International Conference on Hybrid Systems: Computation and Control. pp.297–306 (2016)

#### References

